

# Hybrid Type Systems

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## Type systems

Type systems are a lightweight verification method

- ▶ Common in programming languages
- ▶ Increase software reliability
- ▶ Verify basic interface specifications
- ▶ Avoid complicated formalism

# Type systems

Static  
multiple types

- ▶ Earlier error detection
- ▶ Better documentation
- ▶ Allow more optimizations
- ▶ Increased runtime efficiency

# Type systems

## Dynamic

type Dynamic

- ▶ More expressive
- ▶ Fast adaptation to requirements
- ▶ Simpler component interaction
- ▶ Truly dynamic behavior

# Problem

- ▶ Choosing between static/dynamic is not obvious
- ▶ Stronger formalism  $\Leftrightarrow$  less flexibility

## Hybrid type systems

### Research goal

- ▶ Develop a hybrid type system
- ▶ Combine best of both static/dynamic
- ▶ Adjust type system to the development process

## Type system properties

- ▶ Gradual typing (introduced by Siek [2])
- ▶ Type inference
- ▶ Polymorphism
- ▶ Generics & heterogeneous data structures
- ▶ Specifications
- ▶ Subtyping & covariance
- ▶ ...

## Gradual typing & Type inference

Type annotations are optional and gradually strengthen the type system

```
// accepted  
  (fn (x:Num) => x + 1) 1
```

```
// rejected  
  (fn (x:Num) => x + 1) true
```



## Gradual typing & Type inference

```
// accepted, cast failure at runtime  
  (fn (x) => x + 1) true
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```

```
≈ (fn (x:Dyn) => x + 1) true
```

```
≈ (fn (x:Dyn) => (<Num> x) + 1) (<Dyn> true)
```

## Polymorphism

Identity function

```
let idI = (fun (x:Int) => x)
(idI 1) : Int
```

```
let idD = (fun (x:Double) => x)
(idD 2.0) : Double
```

```
let idIL = (fun (x:Int list) => x)
(idIL [1,2]) : Int list
```

## Polymorphism

Polymorphic identity function

```
let id = (fun (x) => x) : a → a
```

```
(id 1) : Int
```

```
(id 2.0) : Double
```

```
(id [1,2]) : Int list
```

## Generics & heterogeneous data structures

`[1, 2, 3] : Int list`

`[1.0, 2.0, 3.0] : Double list`

`[1, 2.0, "Hi"] : Dyn list`

`[1, 2.0, "Hi"] : Int ∨ Double ∨ String list`

## Specifications

Fibonacci sequence with refinement types  
(introduced by Flanagan [2])

```
let Pos0 = {x:Int | x >= 0}
```

```
let rec fib (n:Pos0):Pos0 =  
  if (n < 2)  
  then 1  
  else ((fib (n - 1)) + (fib (n - 2)))
```

## Specifications

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## Work & Conclusions

- ▶ Bidirectional typechecking with polymorphic types (by Dunfield [1])
- ▶ Dynamic type encoding through union types (e.g., Furr [2])
- ▶ Integrate refinement types (by Flanagan [2])

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